



COMPARISON OF LOW ENERGY CONSUMPTION PROTOCOLS IN WIRELESS SENSOR NETWORK

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ABSTRACT

Wireless Sensor Network is an emanating field for research in several aspects of human beings. WSN is a network composed of nodes without any central controller. Sensor network protocols have high ability to be used in each and every field of life. The working of sensor networks depends entirely on the battery power of sensors which cannot be replaced. In this work, cluster based routing protocol is studied precisely. Further leach protocol is modified as MCHAEEC protocol by introducing two schemes efficient cluster head replacement scheme and dual transmitting power level scheme for amplification of signals. The performance of MCHAEEC protocol is shown in terms of network life time and throughput. The cluster head selection in MCHAEEC is more efficient than other protocols because energy is not wasted to select the new cluster head for new round and in other scheme two levels of power are used for amplification of signals either low or high power. Thus using different energy for different types of communication. MCHAEEC is compared with LEACH and MOD-LEACH protocol in the end. MCHAEEC gives better performance than LEACH and MOD-LEACH in terms of network lifetime.

Keywords: Clusters, MCHAEEC, LEACH, MOD-LEACH, Power level, Threshold etc.

I. INTRODUCTION

WSN is a large network of nodes which are interconnected by a communication network. These nodes are autonomous objects which are randomly deployed in any region according to the requirement. The number of sensor nodes in a network can vary from some hundreds of nodes to thousands depending on the requirement of the region. A sensor network consists of low power nodes and low cost multifunctional nodes. These sensor nodes communicate via wireless medium over certain distances. The sensor nodes consist of sensing, processing, power and communication component. Each and every component is having its own functions to perform. The arrangement of these sensor nodes is defined in two ways one in deterministic way and other self-organizing capability. It means that in deterministic way the routes for data transmission is already defined and sensors are placed manually. Whereas in self-organizing way the nodes are scattered randomly in an ad hoc manner. As the energy consideration is also important for setting up the routes. They can provide high quality and fault tolerant network capability [1] [2], however their battery lifetime is limited. Ranging from medical applications to military applications WSN is used. The numbers of nodes in network are potentially large. One of the main important applications of the WSN is to monitor the data of the environment and then further transmit that data to the central point called sink node. In order to initiate some specific action the data is analyzed. That data is analyzed by the sink node or by the network. The most important criteria in wireless sensor network are that large number of sensors should be used and the energy consumption of these sensors should be low so that energy is not wasted.

The most prominent limitation in Wireless Sensor Networks is energy limitation constraint. The number of nodes deployed in any region are high so the monitoring of such nodes becomes very difficult. Once the network is established the node energy keeps on dissipating whenever the information is received or sent to other node in network. As the sensor nodes are of small size due to which big size batteries cannot be embedded in them and it is very difficult for the people to replace



such small sized batteries .So sensors need efficient mechanism for energy utilization. As sensor networks have dynamic nature so nodes can lose its connection in any network. So each node is tightly power constrained. Therefore the entire network has limited lifetime. In order to enhance the network lifetime of network many routing protocol approaches are used. Thus routing protocol is an important factor affecting the energy consumption of the sensor nodes. One of techniques to enhance network lifetime is clustering. Clustering divides the nodes into separate clusters .In this large data coming from different nodes is compresses by using the techniques of aggregation. Therefore, only those packets are transmitted which consist meaningful information; this approach leads to an efficient solution for energy. Data aggregation can be executed with the help of clustering technique [3] [4]. In any wireless network for efficient performance, the protocols should be efficient. Many protocols have been designed which address the power problem. One of the most prominent algorithms which are designed to address power issue is direct transmission algorithms, cluster based algorithms and hop to hop transmission algorithms. In further sections, the work done on cluster based routing in wireless sensor network is discussed and some modifications are made in existing protocol to increase efficiency. Finally, experiments are made to show the network lifetime and throughput.

II. RELATED WORK

In a survey on routing protocols in WSN, it classifies the routing techniques based on network structure into three categories flat based, location based and hierarchical based. Further also these protocols are classified into QoS based, data centric based, mobility based, query based, negotiation based and multipath based protocols. Manufacturing of small nodes capable of transmitting and processing are available nowadays, but hundreds of nodes are deployed for any application in any region. These nodes deployed have limited power so the power capability of the nodes should not be wasted. It means the power capability should be used in a précised manner so that it is not wasted and the network lifetime is enhanced. Although efficient circuit is also an essential requirement for energy efficiency but protocols also play an important role in energy consideration and power management.

To deal with such energy problems direct transmission approach was discussed [5].In direct transmission the node senses data from the environment and directly transmits it to the base station. This method is very much efficient for the security measures because there is direct transmission but one drawback of this approach is the excessive power consumption of the nodes. As the nodes which are far away from the base station die out quickly because they require extra power consumption. To solve this problem further protocols are classified in two classes taken into consideration the energy aware broadcast/multicast problems. The algorithm is classified in the MEB/MEM (minimum energy broadcast/multicast) problem and MLB/MLM (maximum lifetime broadcast) problem in wireless networks. Two main energy metrics that are considered in this approach are minimizing the total power consumption of all the nodes involved in multicast session and maximize the operation time until the power of the one node battery is depleted in multicast session.

In order to solve the issues in direct transmission another approach least transmission energy came into existence. In this the transmission of the data takes place along the network through multi hop. Both direct transmission and least transmission energy lead to same energy problem but the only difference between two was that in least transmission the nodes which are far away from the base station live longer as compare to the nodes which are closer to the base station. The reason for the nearer nodes to die is because all data is routed to base station. However the transmitting most of the data from each node uses much energy. To address this problem Direct Diffusion method is used which involves the data processing and data dissemination[6] .This concept is used in hierarchical clustering mechanism and thus presented clustering based routing protocol[7].Clustering is one of the efficient mechanisms in energy conservation of wireless sensor network. Clustering is useful in data query and broadcast of message [8][9]. In sensor network nodes establish clusters and each node nominates one node as cluster head. The node which is chosen as a cluster head is having maximum energy. It is the cluster head which collects the data from each cluster, aggregates the received data and then transmits it to the base station. In this way lifetime of the network is maximized .Whereas some clustering protocols work on the fact that any node which is in the network can



be elected as a cluster head. Considering clustering mechanism various protocols are developed and each of the protocol has its own functions [10]. Stable Election Protocol (SEP) gives each node a weighted probability of becoming a cluster head[11].The sensor network architecture of Threshold sensitive Energy Efficient sensor network protocol (TEEN) is based on hierarchical grouping where the nodes which are closer to each other form clusters and process goes on until second level is reached. The Adaptive Threshold Energy Efficient sensor network protocol (APTEEN) architecture is same as TEEN. In APTEEN first the previous data is analyzed, then snapshot view of the network is taken and for a specific period of time event is monitored. Power efficient Gathering in Sensor Information Systems (PEGASIS) a chain is formed among the nodes and so that each node can receive and transmit the data, then the collected data moves from node to node and destination node transmits it to the base station. LEACH is among the most popular protocols for the sensor network. Leach [12] protocol proposes that each node should elect itself a cluster head; therefore first CH selection process is initiated. In Leach protocol each node is assumed to have powerful energy to reach the nearest cluster head. In this nodes transmit to cluster head and cluster head aggregate the data and transmit it to the base station. But using powerful energy for transmission waste energy of the network. Based on LEACH and SEP several numerous protocols are proposed Q-Leach optimize the network life time of homogenous network [13].

III. LEACH PROTOCOL

Leach is one of the prominent proactive sensor network protocols. In leach once the cluster head is chosen , it is responsible for broadcasting message in assigned time schedule .As the cluster head broadcasts a TDMA schedule for the transmission of the data in given order. Each and every node is having its own time slot in the frame during which the transmission takes place. When the last node completes its transmission then the schedule is repeated. Leach is a cluster based protocol in which cluster head rotates randomly to evenly distribute the energy among the nodes in the network [14]. Leach protocol works in various rounds. Each round consists of two phases, set up phase and Steady phase.

1. **Setup Phase:** In this phase each node decides whether or not to become cluster head. It depends on decision made by the node by choosing a random number between 0 and 1[15]. The node which is having maximum energy is chosen as a cluster head. Then the cluster head CH broadcasts its message to all clusters that it has become the new cluster head CH. Then all the nodes have to join the cluster, thus all nodes send join request message to the cluster. The nodes join that cluster which is having signal intensity higher. In this way the each CH know its members. Then CH creates time schedule so that each node know its time slot during which the data is transmitted.

2. **Steady Phase:** In this phase the data is send to the cluster head by their nodes in the allocated time slot. Then cluster head CH aggregates the data and sends it to the base station BS. The nodes turn radio on only when it senses some important information. But the cluster head has to keep its receiver on so that it can receive the data from the nodes. The routing topology of leach is shown below in Figure1.

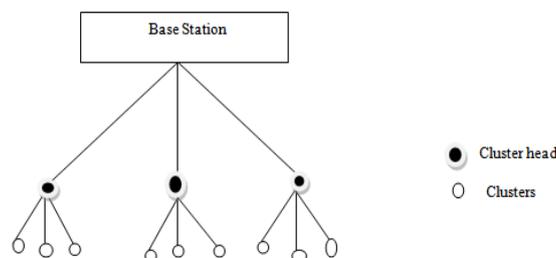


Figure 1 LEACH Routing Topology

The stochastic CH selection algorithm is used in LEACH. Leach is proposed for homogenous network. When CH performs its function of aggregation and transmitting, CH consumes extra energy. Because the node selected as the cluster head dies out quickly due to extra burden. Thus there should be equal distribution of load among all nodes. For this purpose Leach



proposes that every node becomes cluster head CH after $1/p$ rounds [16]. Each sensor node n generates a random number. And this random number is compared to threshold value $T(n)$. If random number is less than threshold value then sensor node becomes cluster head for that particular round. Otherwise becomes cluster member for that round.

$$T(n) = p^{1-p(r \bmod 1/p)} \text{ if } n \in G(1)$$

In this formula p is the probability of the node to be elected as a cluster head, r is the number of rounds selection, G is the set of nodes which are not elected as a cluster head in the given round.

IV. PROPOSED SCHEME

In proposed method two schemes are used Cluster head replacement scheme and two power level scheme for amplification. Basically Leach protocol is modified and in Leach protocol the node elected as cluster head for one round does not become cluster head for other rounds. In this way if the cluster head has some energy left during transmission, that energy is wasted. Because the same node cannot become cluster head in next round. Whereas in the proposed scheme if the cluster head has some energy left during its time period and the energy left is greater than the threshold then same cluster head can be chosen for the next round. In this way the energy is not wasted as in Leach. If the cluster head has less energy than the threshold, then new node can be chosen as a cluster head replacing the old one. In second scheme the two different power levels are introduced for the amplification of signals. In intra cluster communication, transmission takes place within cluster and energy required should be less. Whereas in inter cluster communication, where the transmission takes place within the cluster heads the energy requirement is more. So using same energy for both types of transmission is wastage of energy. In this scheme different power levels are used for the different type of transmissions. In this scheme when node becomes cluster head, routing protocol informs the same node to use the high amplification energy and when the same node becomes the cluster member in another round the protocol informs it to switch it to low power level. Finally soft and hard thresholds are applied on this modified Leach.

The operation of new protocol is divided in two phase's one full transmission and half transmission. The node selected as a cluster head performs the main functions. The remaining nodes in the cluster broadcast their status to the neighbouring nodes. Any node which seems ready to accept this status saves this information and broadcast this message to other nodes that the nodes status indicates that this node can be elected as a new cluster head. In this way confirmation message is sent to the nodes. The nodes which receive the confirmation message send back the acknowledgement to the cluster head with another confirmation of electing that node to become the cluster head. Thus the cluster head saves the id of that node which is approved to become the cluster head for the next round. In this way the old cluster head receives the data, aggregates the data and then send this data to the base station. If the status of the node which is chosen as the next cluster head shows that this node has high energy to be used as a cluster head in next round as well. Then it is called full transmission. Whereas in half transmission, only the sensors with cluster head, will aggregate the data and send the data to their respective cluster head. Then the old cluster head will aggregate the data and transmit it to the base station.

V. PROPOSED PROTOCOL

In this section new protocol MCHAEEC is described.

Cluster head selection: In this modified protocol, in the beginning of each round cluster head is elected. In order to determine the ability of the node to become cluster head, random number is generated between 0 and 1. Then is number is compared with the sensor threshold value $T(n)$. If the threshold value is greater than the random number then the node becomes the cluster head for the current round. This random number is actually a sensor residual energy. This energy is generated by the sensor node to determine the method to check the capability of the node to become cluster head. The threshold value for desired cluster head percentage p , for sensor s in round r , with residual energy R and maximum energy M is calculated as. G is a set of nodes which has not been selected as cluster head in last $1/p$ rounds.



$$T(s) = \left\{ \begin{array}{l} \max \left\{ \frac{p}{1-p(r \bmod 1/p)} \cdot X \cdot R, \frac{RE}{M} \right\} \quad \text{if } s \in G \\ 0 \quad \text{otherwise} \end{array} \right.$$

This is the formula used to calculate the threshold value for sensor node. Based upon the same formula in Mo-Leach the threshold value is calculated for two rounds as follows:

$$T(s) = \left\{ \begin{array}{l} \max \left\{ \frac{p}{1-p(r \bmod 1/p)} \cdot X \cdot \frac{RE}{M^2}, T_{min} \right\} \quad \text{if } s \in G \\ 0 \quad \text{otherwise} \end{array} \right.$$

Where T_{min} is minimum threshold energy, M is maximum energy and RE is residual energy. If the node approve this formula then the node selected as a cluster head is capable to become cluster head for two rounds. Whereas if the same formula is not approved by node then another formula is used for determining the ability of node to become cluster head for one round.

$$T(s) = \left\{ \begin{array}{l} \max \left\{ \frac{p}{1-p(r \bmod 1/p)} \cdot X \cdot \frac{RE \cdot T_{min}}{M}, \frac{RE}{M} \right\} \quad \text{if } s \in G \\ 0 \quad \text{otherwise} \end{array} \right.$$

This is the approach used in selection of the cluster head in order to minimize the energy wastage in selecting the cluster head for next round.

Dual transmitting power levels: In this scheme different power levels are used for different transmissions. Depending on the distance of the nodes, power levels are used for the amplification. Whereas in Leach same amplification energy is used for all transmissions. When the communication is within a cluster, means intra cluster communication. Where the data is sensed by cluster members and reported to cluster head. This is called intra cluster communication. The amplification energy required for intra cluster communication is low. Thus using low power level for intra cluster communication saves energy. When the communication is between two cluster heads, it is called inter cluster communication and there is one more communication that is cluster head to base station communication. The energy requirement cannot be same for all types of communication. Minimum energy is required for inter cluster communication. Whereas energy required for cluster head to base station is high. In other clustering protocols same energy is used for communications. This scheme provides two type of power levels for communication one high power level and low power level. When the node which is used as a cluster head, routing protocol informs it to use high amplification energy. When the same node becomes cluster member in any other round routing protocol switches it to low amplification energy. So different levels of power are used in order to minimize the energy wastage as in other schemes by issuing same energy for different communication.



5.1 Design Diagram

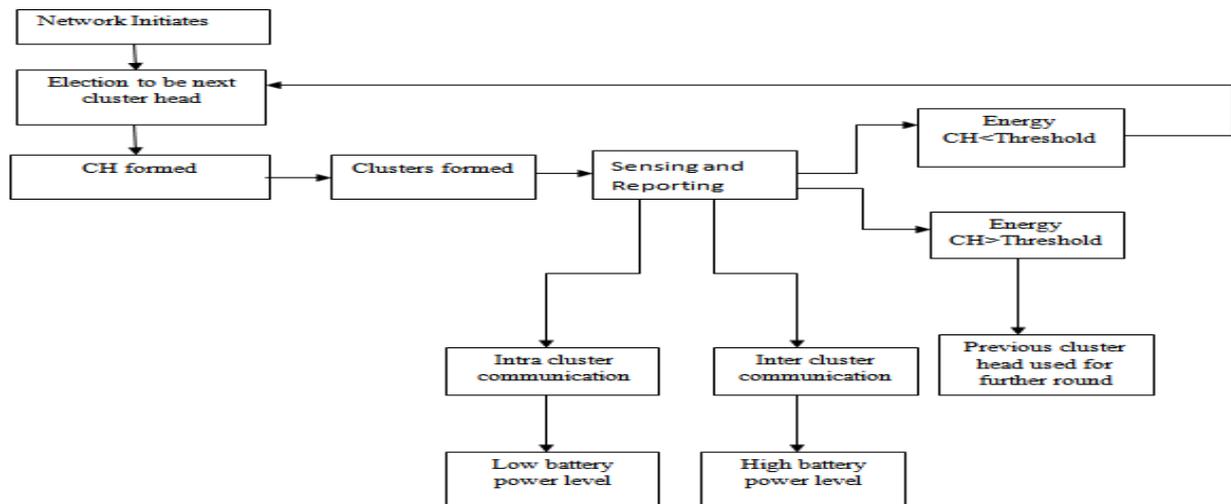


Figure 2 Design Diagram

VI. SIMULATIONS

Simulations are conducted in MATLAB. In order to show the performance of MCHAECC specific intervals and parameters are taken. MCHAECC is compared with LEACH, MOD-LEACH in terms of network lifetime and it is clear from the graph that MCHAECC outperforms LEACH and MOD-LEACH.

6.1 Simulation Scenario and Parameters

In order to perform simulations a network size 100X100 m² is taken. Total nodes taken are 100 with initial energy of sensor nodes 0.5J and packet size 4000 bits. Transceiver idle state energy consumption 50nj/bit with data fusion and aggregation energy consumption 5nj/bit.

6.2 Simulation Results

In a wireless sensor network computing capacity and stored energy is limited. Due to this limiting computing capacities, network lifetime and throughput is affected. Thus in the simulation the ICHEAC is evaluated using following indices.

6.2.1 Network Life Time

The network lifetime of the sensor network is the lifetime of the network from the starting of the network to the end of the network. It means the time from where the network starts its operation till the phase network has completed its operation. The operation is measured in terms of the rounds. Thus the network lifetime is measured in two ways a live nodes and dead nodes.

6.2.2 Throughput

The amount of data received by the base station describes the rate of the accuracy of the nodes, throughput. The more data received means high accuracy. The throughput of the sensor network is measured by the total number of packets sent to base station, packets sent to cluster head during the network lifetime and cluster head formation.

6.3 Simulation Analysis

The implementation of the MCHAECC is shown in form of graphs and the original LEACH and MOD-LEACH protocol is compared to the MCHAECC. Figure.3 and Figure.4 shows the change of a live nodes and dead nodes over the rounds. In Figure 3 it is clear that the number of alive nodes are higher between 0 and 500 rounds. Further as the rounds are increasing ranging from 1000 to 1500, the number of alive nodes are decreasing. In fig 4.the number of dead nodes are shown over the



rounds. It is clear that number of dead nodes are low in initial rounds ranging between 0 and 500. Between 1000 and 1500 rounds the range of dead nodes is higher. The change of alive nodes and dead nodes show the operational time of the MCHAECC protocol.

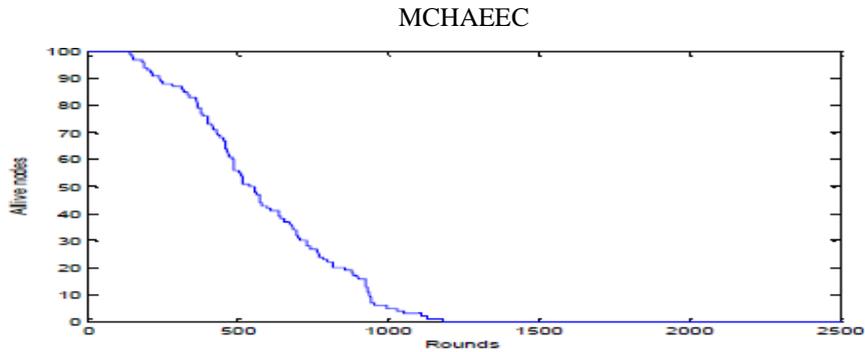


Figure 3 Alive nodes vs. Round

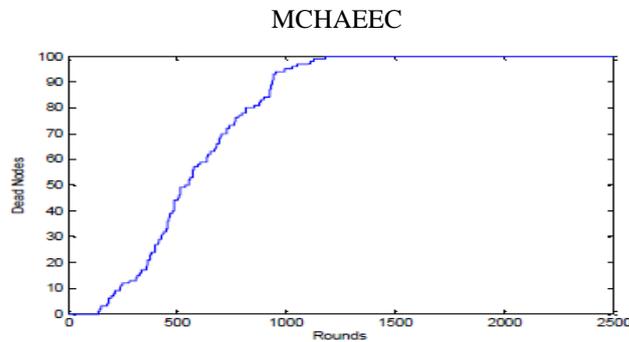


Figure 4 Dead nodes vs. Rounds

In Figure 5 and 6, the analysis of the number of packets sent to Base station and Cluster head is shown over rounds in order to measure accuracy. In Figure 5 it is clear that the packets sent to base station increases as the number of rounds increases. More packets are sent to base station between 1500 and 2500 rounds.

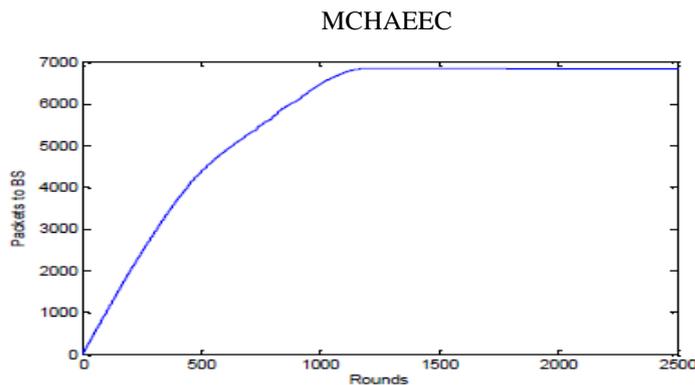


Figure 5 packets to BS between Rounds

MCHAECC

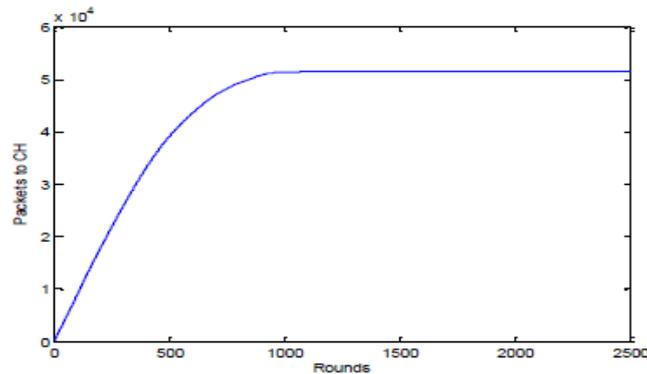


Figure 6 Packets to CH between rounds

In Figure 6 the number of packets sent to the cluster head are shown. The number of packets sent to cluster head also increases as the number of rounds increases. The high number of data packets are sent between 1000 and 2500 rounds. The more the data packets sent to the cluster head the more accuracy of the network lifetime is achieved. This transmission shows the accuracy of the network life time. In Figure 7 the cluster head formation is shown over rounds. The variation of the cluster head formation is shown between 0 and 1000 rounds.

MCHAECC

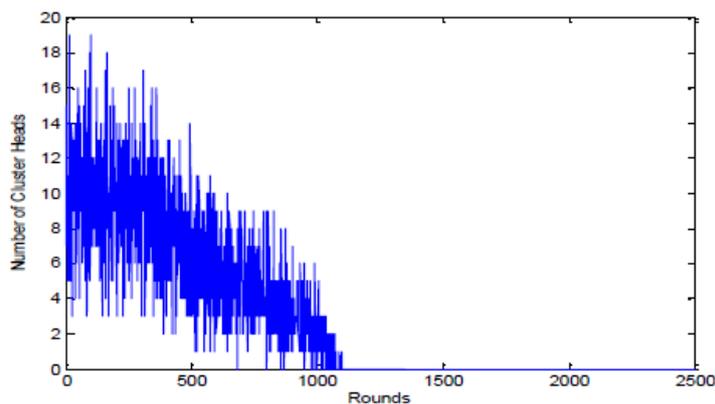


Figure 7 Number of Cluster head formed vs. rounds

In Figure 8 MCHAECC is compared with LEACH and MOD-Leach in terms of network life time and MCHAECC outperforms both the protocols.

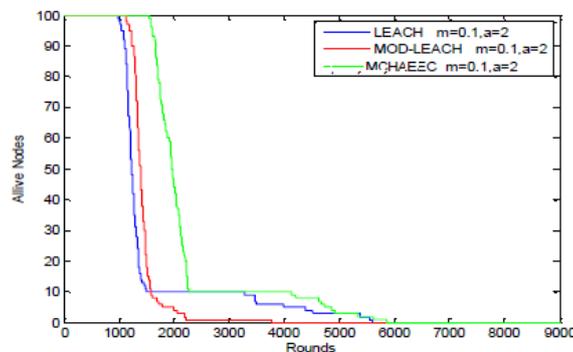


Figure 8 Comparison of LEACH, MOD-LEACH and MCHAECC



VII. CONCLUSION AND FUTURE WORK

In this work new clustering approach is proposed MCHAECC, which can be further, utilized in other clustering routing protocols for better efficiency. In this efficient cluster head replacement scheme is used and dual transmitting power level. Whenever cluster head is selected in first round then the same node can be used as a cluster head in other round if the energy of the node is greater than the threshold. In this way the energy of the cluster head is not wasted. If the energy of the cluster head is less than the threshold then the other node is chosen as a cluster head for next round. In this method the threshold is given priority for the choosing of the cluster head. And for different types of communication different power levels are used either low or high. Depending upon the distance the different amplification energy is used for the different types of communication.

Further the work on data aggregation and fusion can be carried out in cluster selection. Hence there is need to develop the new protocol which spends least energy in receiving and fusion of the data to base station and much more efficient mechanisms for the cluster head replacement.

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