



Non Redundant Radix-4 Signed Digit encoding DSP Accelerator

Subhashini¹, C. Madhavi²

¹M.Tech (DSCE), ²Assistant Professor (ECE), Sri Visvesvaraya Institute of Technology & Science, Chowderpally, Devarkadra, Mahabubnagar

ABSTRACT

Dsp accelerator has very prominent role in digital signal processing applications. Acceleration mainly performed in the application specific integrated circuit devices. Dsp accelerator has perform various arithmetic operations to improve the performance of the acceleration proposed technique is implemented i.e NR4SD.i.e.Non Redundant Radix-4signed digit encoding technique, it has high performance compared with the modified booth algorithm and also it has less decoding time for data transmission.In this paper the complete design of dsp accelerator is designed with NR4SD technique to improve the area, delay and power consumption. It has implemented with the verilog hdl and synthesize and simulated with the XILINX 14.5 ISE standard.

Keywords: NR4SD, Pre encoded multiplier, FCU, carry save form, modified booth encoding.

I. INTRODUCTION

Modern embedded systems target high-end application domains requiring economical implementations of computationally intensive digital signal process (DSP) functions. The incorporation of heterogeneity through specialised hardware accelerators improves performance and reduces energy consumption. Although application-specific integrated circuits (ASICs) kind the perfect acceleration solution in terms of performance and power, their inflexibility leads to magnified atomic number 14 quality, as multiple instantiated ASICs are required to accelerate varied kernels. Several researchers have proposed the utilization of domain-specific coarse-grained reconfigurable accelerators so as to extend ASICs' flexibility while not significantly compromising their performance.

Highly versatile datapaths have been projected to expeditiously map primitive or chained operations found within the initial data-flow graph (DFG) of a kernel. The templates of complicated chained operations are either extracted directly from the kernel's DFG or laid out in a pre-defined activity template library. Selections on the accelerator's datapath highly impact its potency. Existing works on coarse-grained reconfigurable datapaths in the main exploit architecture-level optimizations, e.g., magnified instruction level correspondence (ILP)

The domain specific design generation algorithms vary the type and variety of computation units achieving a custom-made design structure. In, versatile architectures were projected exploiting ILP and operation chaining. Recently, Ansaloni et al. Adopted aggressive operation chaining to change the computation of entire subexpressions mistreatment multiple ALUs with heterogeneous arithmetic options. The said reconfigurable architecture-



seclude arithmetic optimizations throughout the field of study synthesis and consider them solely at the interior circuit structure of primitive components, e.g., adders, through-out the logic synthesis.

However, research activities have shown that the arithmetic optimizations at higher abstraction levels than the structural circuit one considerably impact on the datapath performance. In, timing-driven optimizations supported carry-save (CS) arithmetic were performed at the post-Register Transfer Level (RTL) style stage. In, common subexpression elimination in metal computations is employed to optimize linear DSP circuits. Verma et al. developed transformation techniques on the application's DFG to maximize.

The utilization of metal arithmetic previous the actual datapath synthesis. The said metal improvement approaches target inflexible datapath, i.e., ASIC, implementations. Recently, Xydis et al. projected a versatile design combining the ILP and pipelining techniques with the CS-aware operation chaining. However, all the said solutions feature an inherent limitation, i.e., metal improvement is finite to merging only additions / subtractions. A metal to binary conversion is inserted before every operation that differs from addition/subtraction, e.g., multiplication, thus, allocating multiple metal to binary conversions that heavily degrades performance owing to long carry propagations. In this transient, we have a tendency to propose a superior field of study theme for the synthesis of versatile hardware DSP accelerators by combining optimization techniques from each the design and arithmetic levels of abstraction. we have a tendency to introduce a versatile datapath design that exploits metal optimized templates of enchain operations. The projected architecture includes versatile procedure units (FCUs), which enable the execution of an outsized set of operation templates found in DSP kernels.

II. MODIFIED BOOTH ALGORITHM

Modified booth algorithm is used for signed digit multiplication, both A and B variable consist n=2k bits and B represented as in the modified booth form that is shown in below.

$$B = \langle b_{n-1} \dots b_0 \rangle_{2's} = -b_{2k-1}2^{2k-1} + \sum_{i=0}^{2k-2} b_i 2^i$$

$$= \langle \mathbf{b}_{k-1}^{MB} \dots \mathbf{b}_0^{MB} \rangle_{MB} = \sum_{j=0}^{k-1} \mathbf{b}_j^{MB} 2^{2j}$$

B is formed with in the limits of {-2,-1,0,1,2}, and

$$\mathbf{b}_j^{MB} = -2b_{2j+1} + b_{2j} + b_{2j-1}, \tag{2}$$

So, B is divided into three digit set like $b_{2j-1}, b_{2j}, b_{2j+1}$, it can be denoted as signed bits depends on MSB bits, each set has individual operation that has shown in below table:

TABLE 1
Modified Booth Encoding

b_{2j+1}	b_{2j}	b_{2j-1}	\mathbf{b}_j^{MB}	s_j	one_j	two_j
0	0	0	0	0	0	0
0	0	1	+1	0	1	0
0	1	0	+1	0	1	0
0	1	1	+2	0	0	1
1	0	0	-2	1	0	1
1	0	1	-1	1	1	0
1	1	0	-1	1	1	0
1	1	1	0	1	0	0

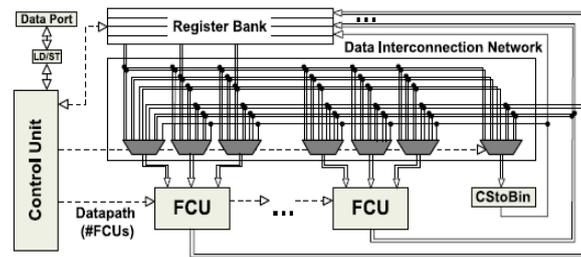


Fig. 1. Abstract form of the flexible datapath.

III. PROPOSED FLEXIBLE ACCELERATOR

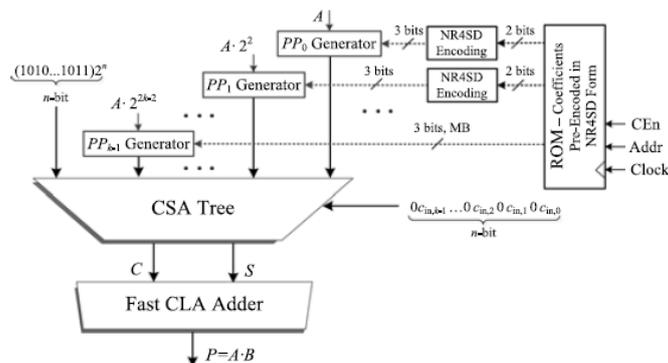
The projected versatile accelerator design is shown in Fig. 1. Each FCU operates directly on metal operands and produces knowledge in the same form for direct use of intermediate results. Each FCU operates on 16-bit operands.

Such a bit-length is adequate for the most DSP datapaths, however the field of study construct of the FCU can be foursquare custom-made for smaller or larger bit-lengths. The number of FCUs is decided at style time supported the ILP and space constraints obligatory by the designer. The CStoBin module could be a ripple-carry adder and converts the metal kind to the two's complement one. The register bank consists of scratch registers and is used for storing intermediate results and sharing operands among the FCUs. completely different DSP kernels (i.e., completely different register allocation and electronic communication patterns per kernel) is mapped onto the projected design mistreatment post-RTL datapath interconnection sharing techniques. The management unit drives the architecture (i.e., communication between the in-fo port and also the register bank, configuration words of the FCUs and choice signals for the multiplexers) in every clock cycle.

IV. NON REDUNDANT RADIX4 SIGNED DIGIT ENCODNG

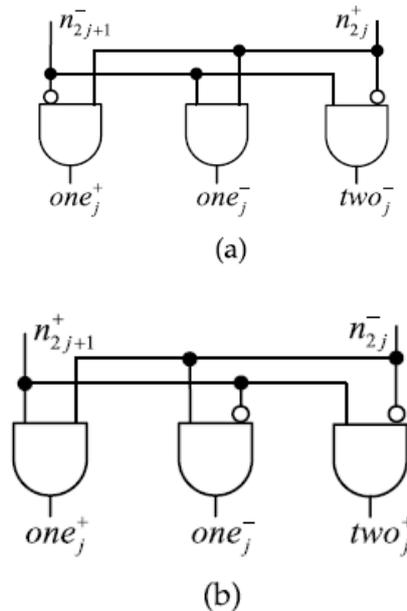
In this section Non redundant radix-4 signed digit encoding has represented, which is advanced method to the modified booth algorithm technique. Modified booth algorithm has worked for the moduli set $\{-2,-1,0,1,2\}$ but this proposed technique has divided into NR4SD+ and NR4SD- with respect to $\{-2,-1,0,1\}$ and $\{-1,0,1,2\}$ moduli sets.

The proposed FCU is designed with the NR4SD techniques to improve the performance. The block diagram of proposed multiplier with NR4SD is shown in the below figure:

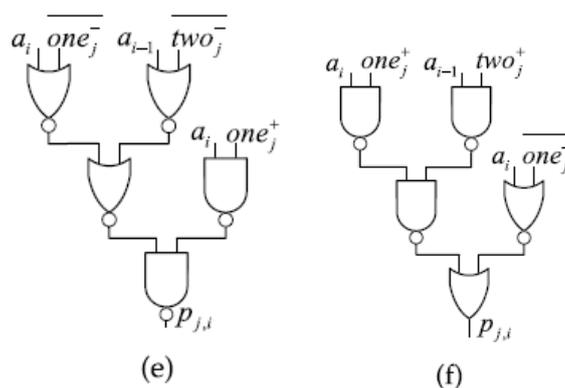


The above diagram represents the proposed multiplier with NR4SD technique, in this the variable B has divided into the sets this are initially stored in the memory ROM. These bits are nothing but n_{2j+1}^-, n_{2j}^+ for NR4SD+ and n_{2j+1}^+, n_{2j}^- for NR4SD-.

These implementations are shown in the below figure.



From the bits of n_{2j+1}^-, n_{2j}^+ and n_{2j+1}^+, n_{2j}^- , three moduli bits are generated $one_j^+, one_j^-, two_j^-$ and $one_j^+, one_j^-, two_j^+$, these are shown in the above diagram. From these the required partial products PP bits are generated as shown in the below figures for NR4SD+ and NR4SD-.



IV. IMPLEMENTATION RESULTS

We executed in Verilog the multiplier plans of Table 5. The PPGs for the NR4SD, NR4SDp multipliers (Figs. 4c and 4d, separately) contain countless since all the A bits are supplemented if there should arise an occurrence of a negative digit. In request to maintain a strategic distance from these inverters and, in this way, diminish the territory/ control/deferral of NR4SD, NR4SDp pre-encoded multipliers, the PPGs for the NR4SD, NR4SDp multipliers were outlined in view of primitive NAND and NOR entryways, and supplanted by Figs. 4e and 4f, separately. The CSA tree and CLA snake were transported in from Synopsys. DesignWare library. The ROM

for the 2's supplement or preencoded coefficients is a synchronous ROM of 512 words regularly met at DSP frameworks, e.g., discourse CoDecs or sound sifting [20]. The width of every ROM relies on upon the multiplier engineering (Table 5). A limited state machine synchronized the information stream and the multiplier operation yet was not considered in the region/control figurings.

V. SYNTHESIS AND SIMULATION RESULTS

In this paper the new design NR4SD technique is used to design the FCU unit for DSP acclerlator and its implemented with the XILINX ISE 14.5 simulation tool and implemented with Verilog HDL. The RTL diagram and simulation results are displayed below.

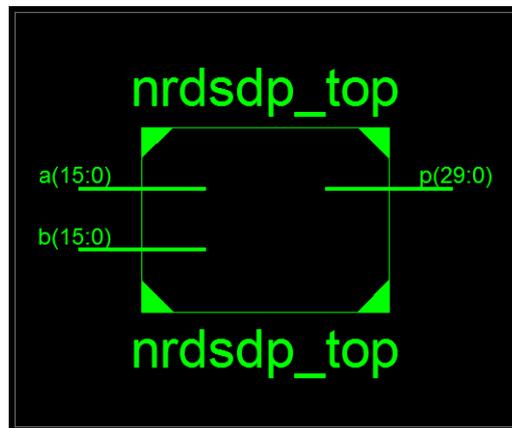


Fig: Top level schematic diagram

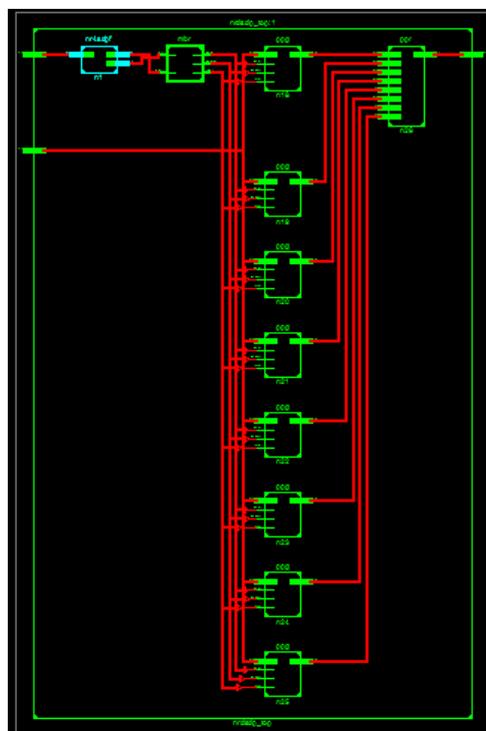


Fig: Internal architectures of RTL diagram

Synthesis result:

Proposed NR4SD based FCU unit results:

Device Utilization Summary (estimated)	
Logic Utilization	Used
Number of Slice LUTs	166
Number of fully used LUT-FF pairs	0
Number of bonded IOBs	158

Total 13.651ns (3.052ns logic, 10.599ns route)
(22.4% logic, 77.6% route)

MB based FCU unit results:

Device Utilization Summary (estimated)	
Logic Utilization	Used
Number of Slice LUTs	311
Number of fully used LUT-FF pairs	0
Number of bonded IOBs	62

Total 14.665ns (2.135ns logic, 12.530ns route)
(14.6% logic, 85.4% route)

Simulation result:



VI. CONCLUSION

In this paper it consist the design of flexible unit with non redundandant radix-4 sign encoding technique for DSP acclearator unit. It performs several arithmetic operations, the experimental results shows the proposed FCU unit has better performance compared with the MB encoding based FCU unit with the area and delay.

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AUTHOR DETAILS

	<p>SUBHASHINI, pursuing M.Tech (DSCE) from Sri Visvesvaraya Institute Of Technology & Science, Chowderpally (Vill), Devarkadra (Mdl), Mahabubnagar (Dist), TS, INDIA.</p>
	<p>C. MADHAVI, working as Assistant professor (ECE) from Sri Visvesvaraya Institute Of Technology & Science, Chowderpally (Vill), Devarkadra (Mdl), Mahabubnagar (Dist), TS, INDIA.</p>