LOW POWER HIGH-PERFORMANCE NON-BINARY LOW-DENSITY PARITY CHECK CODER (FNB-LDPC) OVER GF ($2^m$) USING CHECK-NODE UNIT,

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ABSTRACT
An error correction codes are one of important part in channel coding to enhance the performance of communication system. Recently, low-density parity check (LDPC) codes used as the same purpose but hardware design require large resources, which limits the performance of coders. This paper presents a low power, high-performance non-binary LDPC coder (FNB-LDPC) over Galois fields $GF(2^m)$. The hardware realization of check-node unit (CNU) is a challenging part because it consists of big modules such as FFT/IFFT and multiplier. We overcome these problems by modified CNU structure. The proposed versatile induced flexible LDPC coder supports all field size of suggested Galois fields without the necessity to reconfigure the hardware structure to increase the performance in terms of hardware utilization, power, and delay.

Keywords: FNB-LDPC coder, check node unit, variable node unit, Galois field, real-valued FFT, versatile bit serial multiplier.

I. INTRODUCTION
LDPC codes are a class of block code that satisfies both long length and randomness. The fully parallel LDPC code [1] extended from BP decoding algorithm by capitalizing on parallel structure. A 1024 bit LDPC code achieves a maximum symbol throughput of 1Gbps implemented in ASIC technology. QC-LDPC [2] code achieves high hardware utility efficiency (HUE) and brings about great memory block reduction without any performance degradation. First to split the check matrix into several row blocks, then it performs to improved message passing computations sequentially block by block. A resource efficient LDPC decoder based on a reduced complexity Min-Sum algorithm reduced the interconnect complexity by restricting the extrinsic message length to 2 bits and simplified the CNU [3]. The high throughput decoding of high-rate LDPC codes modified by the sliced message passing (SMP) decoding architecture [4] which overlaps the CNU and VNU and achieved a good tradeoff between area and throughput, and also, high hardware utilization efficiency. A look-up
table (LUT) based VNU design have the best solution for high-speed hardware design and it extended for (2048, 1723) LDPC code of the IEEE 802.3an standard [5]. A parallel NB-LDPC decoder over $GF(256)$ have implemented in 28-nm CMOS technology [6]. The trellis based CNU design maximize the storing capacity by reduced the large amount of memory occupying activities. The sorted log likelihood ratio (LLR) vector of a check-to-variable message has approximated using a piecewise linear function. The first and second minimums have computed by modified CNU [7] in terms of accurate and imprecise manner. In this paper, we present a flexible non-binary LDPC coder (FNB-LDPC) over $GF(2^m)$. The objective of proposed FNB-LDPC coders increases the hardware utilization efficiency by maximum clock frequency, and minimizes the power consumption.

II. PROBLEM DEFINITION AND SYSTEM MODEL
Sulek et al. [18] have proposed a NB-LDPC coder using the mixed domain FFT-BP decoding algorithm with the multiplication units and it also named as semi-parallel decoder. Coder favors mapping a touch of the check to the multiplier focuses embedded in a FPGA, in like way making use of the wide number of sorts of FPGA resources. The throughput wrapped up by a single FPGA by the decoder in light of current conditions made. In NB-LDPC coder, the CNU block enhanced by an approximated evaluation of the nonlinear vectors. The NB-LDPC coder implemented using an FPGA development board from Xilinx with Virtex4 XC4VSX55 device with two different GF orders such as $GF(8)$ and $GF(32)$. The $GF(8)$ NB-LDPC decoder consumes the number of slices utilized as 14535, the number of multiplier as 128 and maximum clock frequency of 170.8MHz. The $GF(32)$NB-LDPC decoder consumes the number of slices utilized as 22494, the number of multiplier as 192 and maximum clock frequency of 130.2MHz. This NB-LDPC coder is not a fully parallel structure and it consume more hardware utilization than existing LDPC coders discussed in related works. Moreover, the multiplier is as main part of CNU block in LDPC coder, but author’s use recursive multiplier for this design. For that reason, we present the flexible non-binary LDPC coder (FNB-LDPC) over $GF(2^m)$ without the necessity of reconfigurable hardware structure. The proposed FNB-LDPC coder implemented over different Galois fields $GF(2^m)$ without modified structure of hardware design. The performance of proposed coder will compare with existing coders including NB-LDPC coder [18].

2.1 System Model of Proposed FNB-LDPC Coder
Check node unit (CNU) module consists of different modules such as permutation, non-linear functions, real-valued FFT/IFFT, and versatile bit serial multiplier.

2.1.1 Real-valued FFT/IFFT module
The initial process in the proposed technique for noise degradation is the transformation of input signals in time domain to the frequency domain. Since speech and noise signals are real valued signals, the conventional FFT
architecture for domain conversion can be replaced with modified low power pipelined architecture so as to make the complete hardware architecture efficient in terms of area and power consumption.

Fig. 1 Parallel Pipelined Architecture for 16 Point Radix 2 RFFT

At stage 1, the butterfly unit will process the pair of real samples $x(\phi)$ and $x(\phi + M/2)$. The butterfly unit consists of 2:1 multiplexer with one selector line $S$. When the inputs are real, then the selector line $S$ set to 1 and the butterfly starts to compute the input values. When the inputs are complex $S$ set to 0, then the multiplexer just passes the input without computation. At stage 2, the architecture consists of shuffling unit, butterfly unit and twiddle factor block A.

The shuffling unit is used to transform the order of the data that required from the stage 1 to stage 2, which also contains 2:1 multiplexer and two delay elements. The twiddle factor ($W^\theta$) module is shown in Fig. 2.

Fig. 2 Twiddle Factor Module
### Table 1: Twiddle factor real and imaginary coefficients for M=16

<table>
<thead>
<tr>
<th>Twiddle factor ((W^\phi))</th>
<th>Real values</th>
<th>Imaginary values</th>
</tr>
</thead>
<tbody>
<tr>
<td>(W^0)</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>(W^1)</td>
<td>0.9239</td>
<td>0.3827</td>
</tr>
<tr>
<td>(W^2)</td>
<td>0.7071</td>
<td>0.7071</td>
</tr>
<tr>
<td>(W^3)</td>
<td>0.3827</td>
<td>0.9239</td>
</tr>
</tbody>
</table>

#### III. EXPERIMENTAL RESULTS

The performance metrics such as device utilization and maximum frequency of proposed and existing coder is given in table 2. The proposed coder implemented with flexible bit design in Virtex7 FPGA devices. The Table 2 shows the performance as device utilization, maximum frequency, and power consumption of proposed router perform very effective than other existing routers.

The fully parallel stochastic LDPC-BC decoder [8] consumes gate count of 760.3K, maximum clock frequency of 768MHz, and power consumption of 437.2mW. The half-stochastic decoding architecture for LDPC-BC decoder over \(GF(16)\) [9] consumes gate count of 1077K gate counts, and maximum clock frequency of 333MHz. The multi-mode LDPC decoder architecture [10] consumes gate count of 320K gate counts, the maximum clock frequency of 400MHz, and power consumption of 284.3mW.

The memory efficient decoder architecture [11] consumes the number of slices is 16803, a number of slice LUT are 31305, the number of slice registers are 4066, the maximum clock frequency of 400MHz, and power consumption of 1638mW.

The Quasi-cyclic LDPC coder [12] consumes the gate counts as 416.2K, maximum clock frequency as 474MHz, and power consumption of 114.3mW. The self-corrected min-sum (SCMS-V1) coder consumes LUT and FF pair count as 60K, maximum clock frequency as 300MHz, and SCMS-V2 consumes LUT and FF pair count as 51K, maximum clock frequency as 300MHz [13].

LDPC-BC decoder [14] consumes an area of 1.79 mm\(^2\), maximum clock frequency of 100MHz, and power consumption of 104mW.

The LDPC decoder [15] consumes an area of 3.11 mm\(^2\), maximum clock frequency of 200MHz, and power consumption of 99.2mW. The LDPC decoder [16] consumes an area of 15.75 mm\(^2\), maximum clock frequency of 100MHz, and power consumption of 800mW. The NB-LDPC decoder [17] consumes a gate count of 564K, the maximum clock frequency of 277MHz, and power consumption of 274mW.

The NB-LDPC coder [18] implemented with two different separate GF orders as \(GF(8)\) and \(GF(32)\). The \(GF(8)\) NB-LDPC decoder consumes the number of slices utilized as 14535, the number of multiplier as 128
and maximum clock frequency of 170.8MHz. The $GF(32)$NB-LDPC decoder consumes the number of slices utilized as 22494, the number of multiplier as 192 and maximum clock frequency of 130.2MHz. The proposed design consumes maximum clock frequency as 506.303MHz, and power consumption of 143mW. The FNB-LDPC 32 bit coder consumes slice registers of 2932, slice LUTs of 3523, number of LUT and FF pairs as 4323 for $y=3$ VNUs, maximum clock frequency of 334.437MHz, and power consumption of 143mW. The maximum frequency of proposed coder increase in terms of 196.43% compare to NB-LDPC coder for $GF(8)$ and $y=8$ and also increase in terms of 156.86% compare to NB-LDPC coder for $GF(32)$ and $y=3$.

<table>
<thead>
<tr>
<th>FPGA family</th>
<th>Maximum Frequency</th>
<th>Power consumption</th>
</tr>
</thead>
<tbody>
<tr>
<td>90 nm CMOS tech.</td>
<td>768 MHz</td>
<td>437.2mW</td>
</tr>
<tr>
<td>90 nm CMOS tech.</td>
<td>333 MHz</td>
<td>-</td>
</tr>
<tr>
<td>65 nm CMOS tech.</td>
<td>400 MHz</td>
<td>284.3mW</td>
</tr>
<tr>
<td>Vitex4</td>
<td>82 MHz</td>
<td>1638mW</td>
</tr>
<tr>
<td>Vitex5</td>
<td>474 MHz</td>
<td>114.3mW</td>
</tr>
<tr>
<td>Vitex7</td>
<td>300 MHz</td>
<td>-</td>
</tr>
<tr>
<td>130 nm CMOS tech.</td>
<td>100 MHz</td>
<td>104mW</td>
</tr>
<tr>
<td>220 nm CMOS tech.</td>
<td>200 MHz</td>
<td>99.2mW</td>
</tr>
<tr>
<td>180 nm CMOS tech.</td>
<td>100 MHz</td>
<td>800mW</td>
</tr>
<tr>
<td>90 nm CMOS tech.</td>
<td>277 MHz</td>
<td>274mW</td>
</tr>
<tr>
<td>Vitex4 (XC4VSX55)</td>
<td>170.8 MHz</td>
<td>-</td>
</tr>
<tr>
<td></td>
<td>130.2 MHz</td>
<td>-</td>
</tr>
<tr>
<td>Vitex7 (XC7VX330T)</td>
<td>317.068MHz</td>
<td>143mW</td>
</tr>
<tr>
<td></td>
<td>334.437MHz</td>
<td>143mW</td>
</tr>
</tbody>
</table>
IV. CONCLUSION

The new flexible non-binary LDPC (FNB-LDPC) coder proposed to enhance the performance of design and it is not required reconfigure hardware structure for any GFs. The proposed coder consists of two units such as check node unit (CNU) and variable node unit (VNU). The hardware realization of check-node unit (CNU) is a challenging part because it consists of big modules such as FFT/IFFT and multiplier.

The proposed FNB-LDPC coder have implemented on Virtex7 FPGA expertise in Xilinx tool. From the experimental results, proposed coder perform very effective than existing coders in terms of hardware utilization, power and delay.

Fig. 3 Maximum Clock Frequency Comparison of Proposed and Existing Work

Fig. 4 Power Consumption Comparison of Proposed and Existing Works
REFERENCES


