

Online Transparent Test for Permanent Faults in First In First Out Buffers

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ABSTRACT

This paper introduces an on-line transparent test procedure for identification of hidden tough errors which improve in firstinput firstouptput buffers of routers through field process of NoC. The method engages replicating tests occasionally to avoid storage of errors. A prototype deign of the recommended test algorithm has been combined into the router-channel barrier and on-line test has been achieved with artificial self-parallel data traffic. The presentation of the NoC following summation of the test circuit has been inspected in requisites of throughput whereas the region transparency has been considered by synthesizing the test hardware. Moreover, an on-line test method for the steering logic has been advised which considers utilizing the title darts of the data transfer progress in carrying the test patterns.

Keywords- FIFO buffers, in-field test, NoC, transparent test.

I. INTRODUCTION

Over the last decade, network-on-chip (NoC) has appeared as a better communication infrastructure compared with bus-based communication network for complex chip designs avoiding the difficulties related to bandwidth, signal integrity, and power dissipation. On the other hand, like all additional systems-on-a-chip (SoCs), NoC-based SoCs must too be tested for faults. Testing the constituents of the NoC infrastructure engages testing routers and interrouter connections.

Major quantity of region of the NoC information transfer medium is engaged by routers, which is largely occupied by FIFO buffers and routing logic. Consequently, the possibilities of run-time mistakes or errors happening in buffers and logic are significantly greater compared with the other components of the NoC. Thus, test method for the NoC transportation must start with test of buffers and routing logic of the routers. Moreover, the test must be presented sporadically to guarantee that no error gets collected.

The sporadic run-time practical mistakes have been solitary of the main apprehensions through testing of extremely balanced CMOS-supported memories. These mistakes are a artifact of substantial consequences, like as ecological weakness, timing, and less supply voltage and thus are *intermittent* (nonpermanent indicating device damage or malfunction) in nature. However, these *intermittent* mistakes regularly reveal a comparatively great incidence rate and finally be likely to turn into permanent. Furthermore, exhaust of memories too origin intermittent faults to become frequent enough to be classified as permanent. Thus, there is a requirement for online

test method that would identify the run-time mistakes, which are broken in environment but regularly turn into stable eventually.

A. Contribution

In this thesis, we have projected an online apparent test procedure for first-input first-output (FIFO) buffers and steering logic there inside the routers of the NoC transportation. Our involvements are as trails. An apparent SOA-MATS++ test generation algorithm has proposed targeting in-field eternal mistakes enhanced in SRAM-based FIFO memories and it has been developed to operate online and sporadic test of FIFO memory there inside the routers of the NoC. Moreover, we have too projected an online test method for the steering logic that is achieved concurrently with the analysis of buffers. The proposal contributes two approaches of consuming the unexploited portion of the title darts of the incoming data packets in transporting the test patterns. First, deterministic test models for the routing logic produced by *Tetramax* are situated in the unexploited fields of the title dart and are moved through the common cycle. Subsequent, the pseudorandom models in the artificial data traffic utilized through usual process and incoming at the steering logic are considered as test patterns. Fault coverage is calculated for each of the two proposals.

B. Fault Models Considered for the Work

The run-time *permanent* faults considered in this brief are assumed to be *intermittent* mistakes, which have turn into eternal eventually. Accordingly, the fault models believed in this thesis are that of intermittent mistakes. The main factors that direct to intermitent mistakes are timing results, for instance time-dependent dielectric breakdown (TDDB), electromigration, negative bias temperature instability (NBTI), and hot carrier injection (HCI). TDDB is a event wherever the oxide below the gate material of an MOSFET humiliates over time resulting in a short circuit, which are sculpted as stuck-at-faults. Electromigration decreases interrelate conductivity with way of time and directs to open circuit. The open circuits sourced by electromigration are formed as stuck-open-faults. NBTI and HCI boost the threshold voltage of transistors directing to reduce in mobility. Consequently, the performances of the memory core reduces bringing in read and write failures. The write failures are sculpted as transition errors, while read failures are sculpted as read disturb faults.

To review, the objective fault models believed for this thesis are stuck-at fault, stuck-open fault, read disturb fault, and transition fault. Exhaustive performance of these faults can be found.

II.RELATED WORK

As error tolerance in NoC scheme has increased consequence among investigate society, many of articles have been distributed enveloping unusual features of fault tolerance, such as failure methods, fault modeling, analysis, and so on. A complete review precisising the research work in these papers has been produced. Above the years, researchers have suggested many of Design- For-Testability (DFT) systems for NoC transportation testing (testing routers as well as NoC interconnect) and for NoC based core testing. Built-in self test (BIST)-based techniques have too been used for checking routers and NoC interconnect. A latest thesis on NoC

and router checking produces a review of the DFT procedures occupied for testing NoC interconnects and routers in meticulous. Additionally to novel test designs, fault tolerant routing algorithms have too been proposed. FIFO buffers in NoC transportation are huge in amount and extend all over the chip. Consequently, possibility of errors is majorly upper for the buffers evaluated with other components of the router. Together online and offline test procedures have been projected for test of FIFO buffers in NoC. The scheme is an offline test procedure (appropriate for the identification of creating error in FIFO buffers) that recommends an allocated BIST controller for FIFO buffers. Online test procedures for the identification of errors in FIFO buffers of NoC routers have been projected. On the other hand, the procedure believes standard cell-based FIFO buffers, as we believe SRAM-based FIFO architectures. As a result, errors considered in this thesis are unusual from those targeted.

To the excellent of our data, no occupation has been accounted in the survey that suggests online test of SRAM-based FIFO buffers there inside routers of NoC transportation. As a result, we inspected online test procedures for SRAM-based FIFOs in common. The analysis exposed that SRAM based FIFOs are checked exploiting either of the trailing two ways, contributed BIST move toward as projected and or allocated BIST proposed. However, together contributed and allocated BIST move towards being offline test procedures not pass to identify permanent errors, which enhance eventually.

III. PROPOSED TRANSPARENT TEST GENERATION

The errors believed in this thesis, if exploited for SRAMs or DRAMs, can be identified employing standard March tests. On the other hand, if the similar set of errors are believed for SRAM-type FIFOs, March test would not be used openly because of the address limitation in SRAM-type FIFOs and as a result we were encouraged to prefer single-order address MATS++ test (SOA-MATS++) for the recognition of mistakes believed in this thesis. The word-oriented SOA-MATS++ test is characterized as $\{ _ (wa); \uparrow (ra,wb); \downarrow (rb,wa); _ (ra) \}$ where, a is the data background and b is the inversion of the data background. \uparrow and \downarrow are increasing and decreasing addressing order of memory, respectively. \downarrow means memory addressing can be increasing or decreasing. Application of SOA-MATS++ test to the FIFO occupies writing patterns into the FIFO memory and reading them back. Consequently, the memory indexes are destroyed. However, online memory test techniques involve the restitution of the memory indexes after test. As a result, researchers have changed the March tests to apparent March test so that tests can be operated without the requirement of external data background and the memory indexes can be renovated after test. We have accordingly converted the SOA-MATS++ test to apparent SOA-MATS++ (TSOA-MATS++) test that can be applied for online test of FIFO buffers. The apparent SOA-MATS++ test produced is characterized as $\{ \uparrow (rx, w^{-x}, r^{-x}, wx, rx) \}$.

The operations performed through the test signify three stages of the test, to be exact, *invert* stage, *restore* stage, and *read* stage. The initial two functions form a read write pair (rx, wx) signifying the invert phase where the initial content (content before start of test) of the FIFO buffer position under test (lut) is read and its inversion is written back to the same position. The invert phase is trailed by reinstat phase engaging the functions (rx, wx) , where the content of lut are read and reinverted. At this point of the test, the contents of lut have flipped double to retrieve the novel content. The last stage, (rx) occupies reading the content of lut exclusive of any write function to trail.

A. Test Algorithm

The algorithmic explanation of the apparent SOA-MATS++ test is illustrated in Algorithm 1. It explains the step-by-step procedure to operate the three stages of the apparent SOA-MATS++ test for every position of the FIFO memory. The objective position for test is specified by the loop index i to fluctuate from 0 to $N - 1$, where N is the apparent SOA-MATS++ Test Algorithm amount of positions in the FIFO memory. In other terms, i denotes the address of the FIFO memory location currently under test. For every position, the three test runs are presented through three stages of the loop index j .

Algorithm 1 Transparent SOA-MATS++ Test Algorithm

```

Require: N = number of rows of the FIFO memory
1:  $i \leftarrow 0$ ; /* memory address pointer */
2: while ( $i \leq N - 1$ ) do
3:    $j \leftarrow 0$ ; /* test cycle counter */
4:   while ( $j \leq 2$ ) do
5:      $temp \leftarrow read(i)$ ;
6:     if ( $j = 0$ ) then
7:        $original \leftarrow temp$ ;
8:        $write(i, !temp)$ ;
9:     else
10:      if ( $j = 1$ ) then
11:         $result \leftarrow compare(temp, original)$ ;
12:         $write(i, !temp)$ ;
13:      end if
14:    else
15:       $result \leftarrow compare(temp, original)$ ;
16:    end if
17:     $j \leftarrow j + 1$ ;
18:  end while
19:   $i \leftarrow i + 1$ ;
20: end while

```

Algorithm 1. Transparent SOA-MATS++ Test Algorithm

For a meticulous FIFO memory position (present value of i), the primary stage of j (address run1) functions the invert phase, where the content of the FIFO position is reversed. The invert test phase occupies reading the content of lut kept on a transitory variable $temp$ and then supporting it up in $original$. Then, the complemented content of $temp$ is written back to lut . At this moment, the content of lut is complement of content of $original$.

In the subsequently stage of j (address run2), the reinstatement phase is operated. The content of lut is reread into $temp$ and evaluated with the content of $original$. The evaluation can outcome in all 1's pattern. However, deviation from the all 1's pattern at any bit position represents fault at that meticulous bit position. Subsequently, the complemented content of $temp$ is written back to lut . As a result, the content of lut , which were inverted after the primary stage get restored after the second.

The third stage of j operates only a read function of lut , wherever the content of lut is read into $temp$ and evaluated with the contents of $original$. At this stage of the test, all 0's pattern in the result indicates fault free location, while deviation at any bit position from all 0's pattern means fault at that meticulous bit location. The final read function ensures the identification of faults, which remained un-

detected during the previous two test runs. At the finish of the three test runs (iterations of j), the loop index i is incremented by one to mark the begin of test for the next location.

B. Fault Coverage of the Proposed Algorithm

For test of stuck-at fault, transition fault, and read disturb fault tests enhanced through field function of FIFO memories is the apparent SOA-MATS++ algorithm. The error exposure of the algorithm is illustrated in Fig. 1. In mutually the figures, the word size of FIFO memory is supposed to be of 4 bits. The text in italics against the arrows denotes the operation functioned, while the text in bold font communicates to the variables used in Algorithm 1.

As illustrated in Fig. 1, assume the data word present in *lut* be 1010. The test cycles start with the reverse stage (memory address pointer j with 0 value) through which the content of location addressed is read into *temp* and then backed up in the *original*. The data written back to *lut* is the complement of content of *temp*. Thus, by the finish of the cycle, the data perform in *temp* and *original* is 1010, while *lut* includes 0101. Let a stuck-at-1 error at the most significant bit (MSB) position of the term accumulated in *lut*. Therefore, in place of accumulating 0101, it essentially accumulates 1101 and accordingly, the stuck-at-error at the MSB obtains motivated.

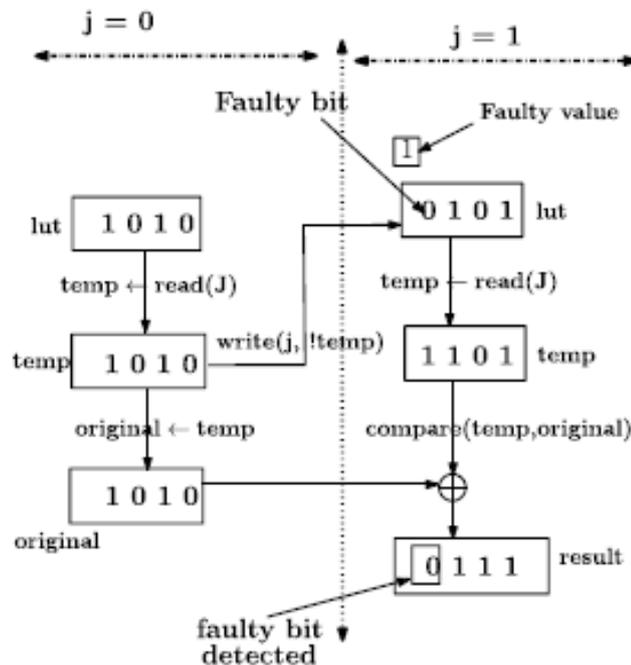


Fig. 1. Fault detection during invert phase and restore phase of the transparent SOA-MATS++ test

During the next stage of j , when *lut* is readdressed, the data read into *temp* is 1101. At this moment, the data perform in *temp* and *original* are compared (bitwise XORed). An all 1's pattern is projected as result. Any 0 within the pattern would mean a stuck-at fault at that bit position. This situation is illustrated in Fig. 1, where the XOR of 1010 and 1101 yields a 0 at the MSB location of the *result* indicating a stuck-at-error at the MSB position. On the other hand, for cases wherever the primary data

for a bit location is unusual from the faulty bit value, the stuck-at-fault cannot be detected for the bit position after the reinstate phase of the test. It needs another test cycle to excite such faults.

IV. IMPLEMENTATION OF THE TEST ON FIFO BUFFERS OF NOC ROUTERS

In this section, we present the method used for performing the projected apparent SOA-MATS++ test on a mesh-type NoC. Data packets are separated into flow control units (*flits*) and are sent in pipeline fashion. The data progress in a mesh-type NoC infrastructure believed for this work is assumed to require buffering only at the input channels of routers. Thus, for a data traffic progress from single core to different, the online test is operated only on the eexcite channel FIFO buffers, which recline beside the path. The buffers perform in two modes, the *normal* mode and the *test* mode. The normal mode and test mode of operation of a FIFO buffer are coordinated with two different clocks. The clock used for test intention (referred as *test_clk* in this thesis) is an earlier clock evaluated with the clock required for normal mode (*router clock*).

The FIFO buffers are allowed to be effective in normal mode for adequate quantity of time before starting their test procedure. This delay in test start produces adequate time for run-time intermittent faults enhanced in FIFO buffers to convert into eternal faults. The test procedure of a objected FIFO buffer is started by a counter, which controls the FIFO buffer from normal mode to test mode. The switching of FIFO buffers from normal mode to test mode happens following a definite period of time without caring about the present state of the FIFO buffer. It can be disputed that at the instant of switching, the buffer may not be full, and accordingly not all locations would be ensured through the test cycle. However, test beginning after the buffer obtains full can cause the following problems. First, wait for the buffer to get complete would unreasonably delay the test initiation process and would allow faults to get accumulated. Second, test of the complete buffer can extend the test time and can negatively affect the normal mode of operation.

A test burst occupies series of test read and write cycles. It needs three read and two write cycles, or in other terms three cycles of the faster test clock to operate an apparent SOA-MATS++ test on a distinct position of a FIFO buffer. It can be disputed that through a test burst, not all FIFO buffer locations are tested or a test of a location can get interrupted. These two problems can be avoided by periodically testing the FIFO buffers. Sporadic testing of a FIFO buffer permits test of a different set of positions of the FIFO buffer in each test burst. Each time the buffer is switched to test mode, the normal process gets interrupted. The FIFO memory position currently addressed in regular mode, at the instant of switching, becomes the objective position for test. Because regular function is suspended at different instants in different test bursts, the locations tested in each burst would be different. As a result, recurring the test bursts for a numeral times on a FIFO buffer would cover the test of each location as the figure of positions in a FIFO buffer is few. Moreover, periodic testing avoids accumulation of error in the buffer.

A. Test Architecture

The FIFO buffer perform in every input channel of an NoC router contains of a SRAM-based FIFO memory of certain depth. Through common function, data darts enter into a *data_in* line of the buffer and are then accumulated in unusual positions of the FIFO memory. On appeal by the neighboring router, the data darts accumulated are exceeded on to the response port through the *data_out* line. Fig. 2(a) represents the FIFO memory with *data_in* and *data_out* line. To perform the apparent SOA-MATS++ test on the FIFO buffer, we inserted a test circuit, few multiplexers and logic gates to the surviving hardware, as illustrated in Fig. 2(a). The read and write functions on the FIFO buffer are controlled by the read enable and write enable lines, correspondingly. The multiplexers *mu6* and *mu7* choose the read and write enable through the normal and test process. Through normal operation when the *test_ctrl* is declared low, the inner write and read enable lines, *wen_int* and *ren_int*, coordinated with the router clock, produce the write and the read enable, respectively. On the other hand, through test

procedure, the write enable and read enable are coordinated with the test clock. As mentioned previous, the read and write functions during test are operated at alternate edges of a test clock. The read functions are coordinated with the positive edges, whereas the *write_clk* is achieved by inverting the test clock. In test mode (*test_ctrl* high), the test read and write addresses are produced by test address generators employed utilizing gray code counters related to the normal address generation. Multiplexers *m4* and *m5* are exploited to select between normal addresses and test addresses.

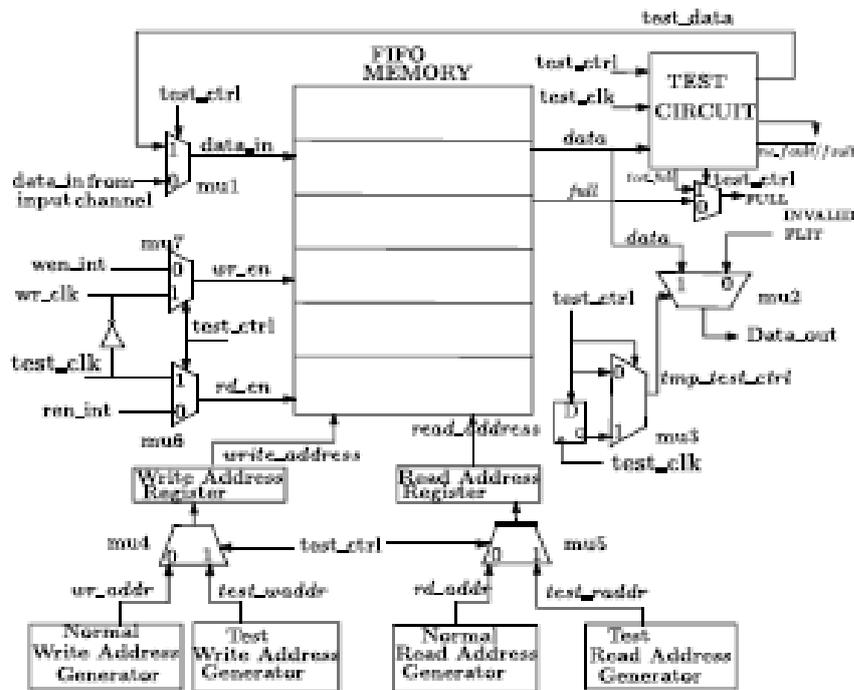


Fig. 2. (a) Hardware implementation of the test process for the FIFO buffers

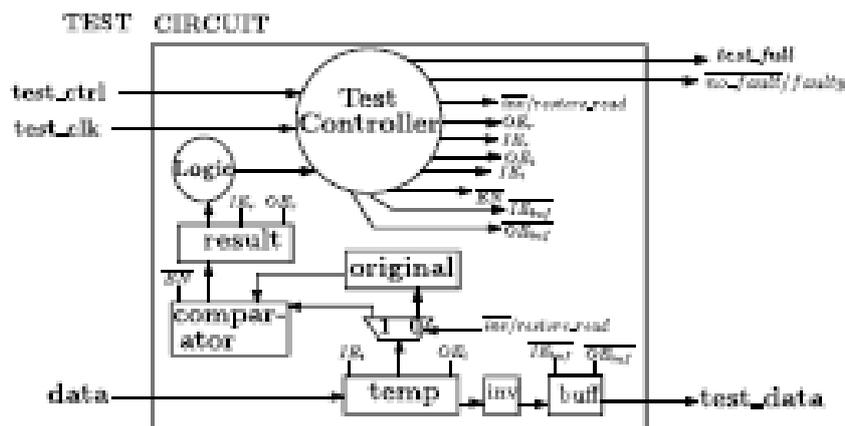


Fig. 2. (b) Implementation of test circuit

Believe the condition while the FIFO buffer is in normal mode with darts being transferred from the memory to the *data_out* line. Later than a some normal cycles, the *test_ctrl* is stated high, controlling the buffer to test mode. Given that the buffer is in test mode, no exterior data is tolerated to be written

to the buffer, or in other terms, the buffer is bolted for the test time. Accordingly, the excite data line for the FIFO memory is changed from the external *data_in* line to *test_data* line presented from the test trail. At the changing moment, the dart which was in the procedure of being moved to the *data_out* line is concurrently read into the *Test Circuit*. However, a single clock cycle delay is produced for the dart to shift to the *data_out* line. This delay guarantees that the dart is not vanished through the changing moment and is accurately received by the router, which requests for it. The distinct cycle delay in the trail of the moving dart is assembled by the D-type flip-flop and the multiplexer *m3*, as illustrated in Fig. 2(a). The dart, which is read in the test trail, is accumulated in a temporary register *temp* and the test procedure begins with this flit.

To avoid packet failure through testing, the *FULL* signal of the FIFO is asserted *high* so that neighboring routers can be avoided from transferring packets to the corresponding router. However, pertaining such procedure increases the network latency as reflected in the results.

V. SYNTHESIS AND SIMULATION RESULTS

The projected FIFO buffers is designed with the XILINX ISE 14.5 simulation tool and implemented with Verilog HDL. The RTL diagram and simulation outcomes are displayed below.

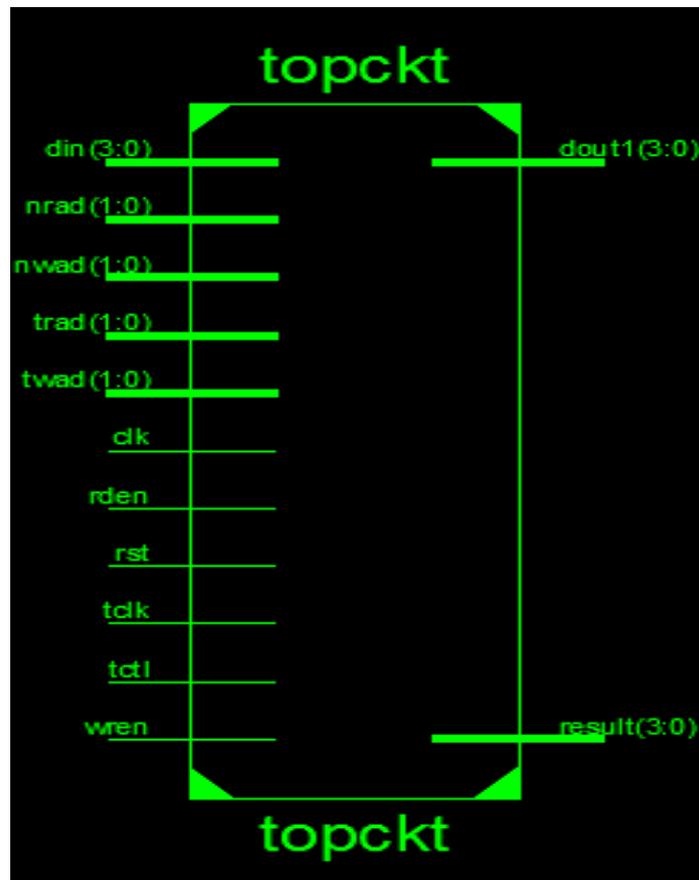


Fig: Top level schematic diagram

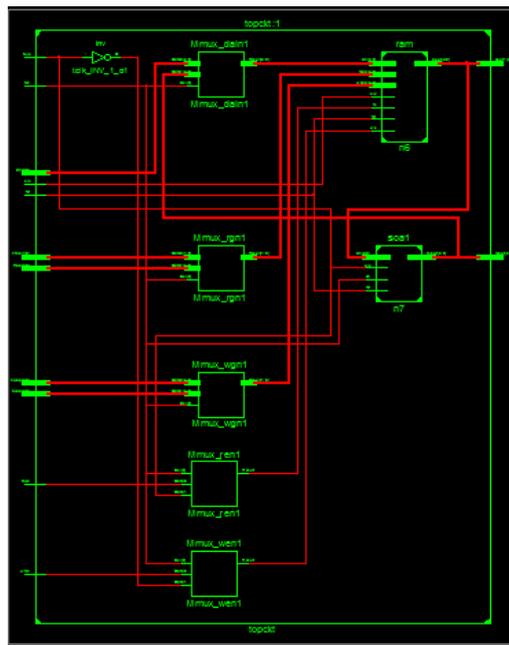


Fig: Internal architectures of RTL diagram

topckt Project Status			
Project File:	FIFOBUFFER..xise	Parser Errors:	No Errors
Module Name:	topckt	Implementation State:	Synthesized
Target Device:	xc7z010-2clg400	Errors:	No Errors
Product Version:	ISE 14.5	Warnings:	No Warnings
Design Goal:	Balanced	Routing Results:	
Design Strategy:	Xilinx Default (unlocked)	Timing Constraints:	
Environment:	System Settings	Final Timing Score:	

Device Utilization Summary (estimated values)				
Logic Utilization	Used	Available	Utilization	
Number of Slice Registers	19	35200	0%	
Number of Slice LUTs	44	17600	0%	
Number of fully used LUT-FF pairs	19	44	43%	
Number of bonded IOBs	26	100	26%	
Number of BUFGB/BUFGCTRL/BUFGCEs	2	80	2%	

Fig: Synthesis report

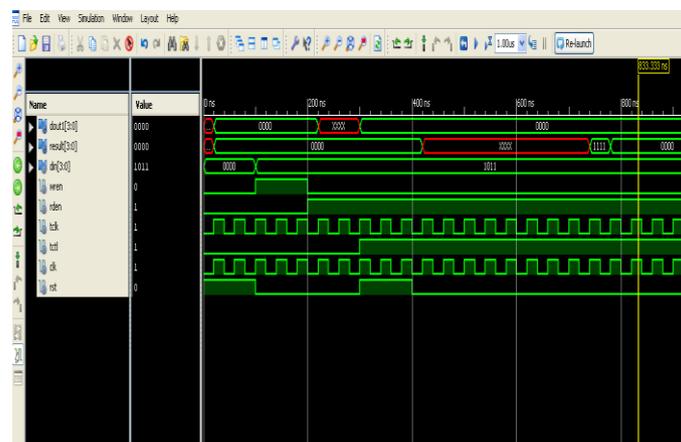


Fig: Simulation result

VI. CONCLUSION

In this thesis, we have proposed transparent SOA-MATS++test generation algorithm that would identify run-time eternal errors grown in SRAM-supported FIFO memories. The projected apparent test is consumed to achieve online and periodic test of FIFO memory present inside the routers of the NoC. Sporadic testing of buffers removes gathering of errors and also tolerates test of every position of the buffer. Simulation results illustrate that sporadic testing of FIFO buffers do not have greatly impact on the general throughput of the NoC excluding when buffers are tested also frequently. We have too projected an online test procedure for the routing logic that is achieved concurrently with the test of buffers and involves utilization of the unexploited fields of the description darts of the received data packets. Test algorithm and RAM was designed by the Verilog HDL synthesized in Xilinx ISE14.5.

VII. FUTURE SCOPE

We have too projected an online test procedure for the routing logic that is achieved concurrently with the test of buffers and involves utilization of the unexploited fields of the description darts of the received data packets. As prospect effort, we would like to change the projected FIFO testing procedure that will tolerate received data packets to the router under test without interrupting the test.

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